

Towards KoS/TTR-based proof-theoretic dialogue management

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Conversational relevance

(a) nice

A: What would you like to cook today?

U: Miso soup

A: How many portions?

U: Four

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(b) not very nice

A: What would you like to cook today?

U: Four portions of miso soup

A: How many portions?

U: ???

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Conversational relevance is crucial for dialogue management. In this respect one should follow dialogue theories (in our case KoS Ginzburg, 2012). There hasn't been much progress in this respect since the early 2000s (Larsson, 2002).

Why (constructive) TTR

1. We use KoS (Ginzburg, 2012) which provides one of the most detailed analysis of conversational relevance
2. KoS is based on the Type Theory with Records (TTR) (Cooper, 2005)
3. Constructive type theories are used as a framework for semantics (e.g. Ranta, 1994)
4. TTR is set-theoretic, and hereby we are presenting proof-theoretic (constructive) version of it; our implementation follows MiniTT (Coquand et al., 2009)

Outline

- **TTR, subtyping and update rules**
- **Dialogue management**
- **Question-answer relevance**
- **Concluding remarks**

Next

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TTR: types

- judgements

$a : Ind$

a is a witness of (the type) Ind (ividual)

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- **ptypes**: $greet(\text{speaker}, \text{addressee})$
- **record types**

$$\left[\begin{array}{l} \text{speaker} : Ind \\ \text{addressee} : Ind \\ \text{content} : greet(\text{speaker}, \text{addressee}) \end{array} \right]$$

TTR: types

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- **ptypes**: $greet(\text{speaker}, \text{addressee})$
- **record types**

$$\left[\begin{array}{l} \text{speaker} : Ind \\ \text{addressee} : Ind \\ \text{content} : greet(\text{speaker}, \text{addressee}) \end{array} \right]$$

- **records**

(1) is a witness of the type above iff $a : Ind$, $b : Ind$, $\theta : greet(\text{speaker}, \text{addressee})$ and the weather doesn't matter.

$$\left[\begin{array}{l} \text{speaker} = a \\ \text{addressee} = b \\ \text{content} = \theta \\ \text{weather} = \text{sunny} \end{array} \right] \quad (1)$$

TTR: type construction operators

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 $(T_1 \rightarrow T_2)$

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- **Meet types** (merge “ \wedge ” in Cooper’s TTR)

$[f : \text{Pasta}] \wedge \begin{bmatrix} f : \text{Cheese} \\ d : \text{Wine} \end{bmatrix}$ reduces to $\begin{bmatrix} f : \text{Pasta} \wedge \text{Cheese} \\ d : \text{Wine} \end{bmatrix}$

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- **Singleton types**

if T is a type and $x:T$, then T_x is a type. $a:T_x$ iff $a = x$.

In record types we use manifest field notation to a represent singleton type. Notations $[a : T_x]$ and $[a = x : T]$ represent the same object.

Subtyping

Assuming that $\text{Pasta} \sqsubseteq \text{Food}$ and $\text{spaghetti} : \text{Pasta}$:

$$[x : \text{Pasta}] \sqsubseteq [x : \text{Food}] \quad (2)$$

$$\left[\begin{array}{l} x : \text{Pasta} \\ c : \text{tasty}(x) \end{array} \right] \sqsubseteq [x : \text{Pasta}] \quad (3)$$

$$[x : \text{Pasta}] \not\sqsubseteq \left[\begin{array}{l} x : \text{Pasta} \\ c : \text{tasty}(x) \end{array} \right] \quad (4)$$

$$[x = \text{spaghetti} : \text{Food}] \sqsubseteq [x : \text{Pasta}] \quad (5)$$

$$[x : \text{Food}_{\text{spaghetti}}] \sqsubseteq [x : \text{Pasta}] \quad (6)$$

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$$[x : \text{Food}_{\text{spaghetti}}] \sqsubseteq [x : \text{Pasta}] \quad (6)$$

Given that $s : S$, $s = [x = \text{spaghetti}]$ and $S = [x : \text{Food}]$

$$S_s \sqsubseteq [x = \text{spaghetti} : \text{Food}] \quad (7)$$

Update rules

1. We maintain dialogue state as a pair of a value and a type

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2. Update rules are defined

$$r : A \rightarrow B$$

3. Applicability condition

$$S_s \sqsubseteq A$$

4. After applying of the rule r state becomes

$$(r(s), B)$$

Example (state and rules)

$$(s_i : S_i), s_1 = [x = \text{spaghetti}], S_1 = [x : \text{Pasta}]$$

$$r_{\text{cook}} : [x : \text{Pasta}] \rightarrow \left[\begin{array}{l} x : \text{Pasta} \\ c : \text{cooked}(x) \end{array} \right]$$

$$r_{\text{cook}} = \lambda s. \left[\begin{array}{l} x = s.x \\ c = \theta_{c(x)} \end{array} \right],$$

where $\theta_{c(x)}$ is a proof that the spaghetti is cooked.

$$r_{\text{serve}} : \left[\begin{array}{l} x : \text{Pasta} \\ c : \text{cooked}(x) \end{array} \right] \rightarrow \left[\begin{array}{l} x : \text{Pasta} \\ c : \text{cooked}(x) \\ d : \text{served}(x) \end{array} \right]$$

$$r_{\text{serve}} = \lambda s. \left[\begin{array}{l} x = s.x \\ c = s.c \\ d = \theta_{d(x)} \end{array} \right],$$

where $\theta_{d(x)}$ is a proof that the spaghetti is served.

Example (change of state)

$$S_1 = [x : \text{Pasta}]$$

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$$S_2 = \left[\begin{array}{l} x : \text{Pasta} \\ c : \text{cooked}(x) \end{array} \right]$$

$$s_2 = r_{\text{cook}}(s_1) = \left[\begin{array}{l} x = \text{spaghetti} \\ c = \theta_{c(x)} \end{array} \right]$$

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$$s_2 = r_{\text{cook}}(s_1) = \begin{bmatrix} x = \text{spaghetti} \\ c = \theta_{c(x)} \end{bmatrix}$$

$$S_3 = \begin{bmatrix} x : \text{Pasta} \\ c : \text{cooked}(x) \\ d : \text{served}(x) \end{bmatrix}$$

$$s_3 = r_{\text{serve}}(s_2) = \begin{bmatrix} x = \text{spaghetti} \\ c = \theta_{c(x)} \\ d = \theta_{d(x)} \end{bmatrix}$$

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Minimal example

Proof-of-concept: basic interaction

U: hello

A: Hello world!

Primary KoS types

$$\textit{InformationState} =_{\text{def}} \left[\begin{array}{l} \text{private} : \left[\text{agenda} : [\text{Move}] \right] \\ \text{dgb} \quad : \left[\text{moves} : [\text{Move}] \right] \end{array} \right]$$
$$\textit{GreetingMove} =_{\text{def}} \left[\begin{array}{l} \text{spkr} \quad : \text{Ind} \\ \text{addr} \quad : \text{Ind} \\ \text{content} : \left[\text{c} : \text{greet}(\text{spkr}, \text{addr}) \right] \end{array} \right]$$

$(\textit{GreetingMove} \sqsubseteq \textit{Move})$

Counter-greet the user

$CGU : InformationState \wedge [\langle LM \text{ is user-greeting} \rangle]$

$\rightarrow InformationState \wedge [\langle agenda \text{ is non-empty} \rangle]$

$CGU = \lambda s. \left[\begin{array}{l} \text{private} = [\text{agenda} = \text{cons}(CGM, s.\text{private}.\text{agenda})] \\ \text{dgb} = s.\text{dgb} \end{array} \right],$

where $CGM = \left[\begin{array}{l} \text{spkr} = \text{agent} \\ \text{addr} = \text{user}_0 \\ \text{content} = [c = \theta] \end{array} \right]$ (θ is a proof of greeting)

Also in the paper

1. Other update rules for integrating user moves and fulfilling the agenda
2. Full example of update chain

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Small extension to TTR

1. Support for boolean types: (*true* : Bool) and (*false* : Bool)
2. Support for conditionals:
 - (IF *true* THEN *x* ELSE *y*) = *x*
 - (IF *false* THEN *x* ELSE *y*) = *y*

Questions

Question definition is a way to establish connection between possible answer and its interpretation in the context of question.

Question : *Type*

Question =_{def} $\left[\begin{array}{l} A : \text{Type} \\ Q : A \rightarrow \text{Prop} \end{array} \right]$

PolarQuestion =_{def} $\left[\begin{array}{l} A = \text{Bool} : \text{Type} \\ Q \quad \quad : A \rightarrow \text{Prop} \end{array} \right]$

UnaryWhQuestion =_{def} $\left[\begin{array}{l} A = \text{Ind} : \text{Type} \\ Q \quad \quad : A \rightarrow \text{Prop} \end{array} \right]$

Example: wh-question

Question : *Type*

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$\llbracket \text{Where do you live?} \rrbracket = \left[\begin{array}{l} A = \text{City} \\ Q = \lambda a. \text{live}(a) \end{array} \right]$

Example: yn-question

Question : *Type*

$$\textit{Question} =_{\text{def}} \left[\begin{array}{l} A : \text{Type} \\ Q : A \rightarrow \text{Prop} \end{array} \right]$$
$$\textit{PolarQuestion} =_{\text{def}} \left[\begin{array}{l} A = \text{Bool} : \text{Type} \\ Q \quad \quad : A \rightarrow \text{Prop} \end{array} \right]$$

[[Do you live in Aix?]] =

$$\left[\begin{array}{l} A = \text{Bool} \\ Q = \lambda a. \text{IF } a \text{ THEN live(Aix) ELSE } \neg\text{live(Aix)} \end{array} \right]$$

Answers

Recalling the definition of question:

$$\textit{Question} : \textit{Type}$$
$$\textit{Question} =_{\textit{def}} \left[\begin{array}{l} A : \textit{Type} \\ Q : A \rightarrow \textit{Prop} \end{array} \right]$$

Answer will return record type for given question:

$$\textit{Answer} : \textit{Question} \rightarrow \textit{Type}$$
$$\textit{Answer} =_{\textit{def}} \lambda q. \left[\begin{array}{ll} \textit{answer} : q.A & \\ \textit{sit} & : q.Q(\textit{answer}) \end{array} \right]$$

Note: $(q.A:\textit{Type})$ and $(q.Q(\textit{answer}):\textit{Prop})$

Example: answer to wh-question

Answer : *Question* \rightarrow *Type*

Answer =_{def} $\lambda q.$ $\left[\begin{array}{l} \text{answer} : q.A \\ \text{sit} \quad : q.Q(\text{answer}) \end{array} \right]$

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$$\llbracket \textit{Where do you live?} \rrbracket = \left[\begin{array}{l} A = \textit{City} \\ Q = \lambda a. \textit{live}(a) \end{array} \right]$$

$$\textit{Answer}(\llbracket \textit{Where do you live?} \rrbracket) = \left[\begin{array}{ll} \textit{answer} : \textit{City} & \\ \textit{sit} & : \textit{live}(\textit{answer}) \end{array} \right]$$

$$\llbracket \textit{in Aix} \rrbracket (\llbracket \textit{Where do you live?} \rrbracket) = \left[\begin{array}{ll} \textit{answer} = \textit{Aix} & \\ \textit{sit} & = \theta_{\textit{la}} \end{array} \right]$$

Example: answer to yn-question

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$\llbracket \text{Do you live in Aix?} \rrbracket =$

$\left[\begin{array}{l} A = \text{Bool} \\ Q = \lambda a. \text{IF } a \text{ THEN live(Aix) ELSE } \neg \text{live(Aix)} \end{array} \right]$

Answer($\llbracket \text{Do you live in Aix?} \rrbracket$) =

$\left[\begin{array}{l} \text{answer} = \text{Bool} \\ \text{sit} \quad = \text{IF answer THEN live(Aix) ELSE } \neg \text{live(Aix)} \end{array} \right]$

$\llbracket \text{yes} \rrbracket(\llbracket \text{Do you live in Aix?} \rrbracket) = \left[\begin{array}{l} \text{answer} = \text{true} \\ \text{sit} \quad = \theta_{\text{la}} \end{array} \right]$

Also in the paper:

1. Dealing with answers in form of propositions (“I live in Aix” instead of “in Aix”)
2. Partial (and incremental) resolution of answers, for utterances like:

A: What do you want today?

U: A beer, please, and chips.

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Concluding remarks

1. Our aim is to maintain tight connection between dialogue theory and dialogue systems' design
2. Dream system:
 - maintains rich information state
 - uses domain-dependent and domain-independent dialogue rules
 - learns probabilities for the rules

Questions?

[speaker=Vlad : Ind
attendees : [Ind]
content : thank(speaker, attendees)]

References I

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- Coquand, T., Kinoshita, Y., Nordström, B., and Takeyama, M. (2009). A simple type-theoretic language: Mini-TT.
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Update rules (Robin Cooper)

Update episode:

$$\lambda r : [\text{agenda} = [[e : \text{pick_up}(a, c)]] : [\text{RecType}]]$$
$$\lambda e : [e : \text{pick_up}(a, c)] .$$
$$[\text{agenda} = [e : \text{attract_attention}(a, b)] : [\text{RecType}]]$$

has a type: $[\text{agenda} : [\text{RecType}]] \rightarrow (\text{Rec} \rightarrow \text{RecType})$

... [it maps] an information state containing an agenda (modelled as a record containing an agenda field) and an event (modelled as a record) to a record type.

Update rules (our approach)

$$IS = [\text{agenda} : [\text{Event}]]$$

$$UR : IS \rightarrow \text{Event} \rightarrow IS$$

$$UR_k : IS \wedge [\text{agenda} = [[e : \text{pick_up}(a, c)]]] : [\text{Event}]$$

$$\rightarrow [e : \text{pick_up}(a, c)]$$

$$\rightarrow IS \wedge [\text{agenda} = [[e : \text{attract_attention}(a, b)]]] : [\text{Event}]$$

and we also specify how the value is assigned:

$$UR_k : \lambda s. \lambda e. [\text{agenda} = [[e = \theta_{(a,b)}]]],$$

where $\theta_{(a,b)} : \text{attract_attention}(a, b)$ is a proof that an event can be counted as such.